

# Verifying Incomplete and Evolving Specifications



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# Verifying Specifications

Requirements  $\phi$



Specification M



Verification



yes



no +

counterexample



# Evolving

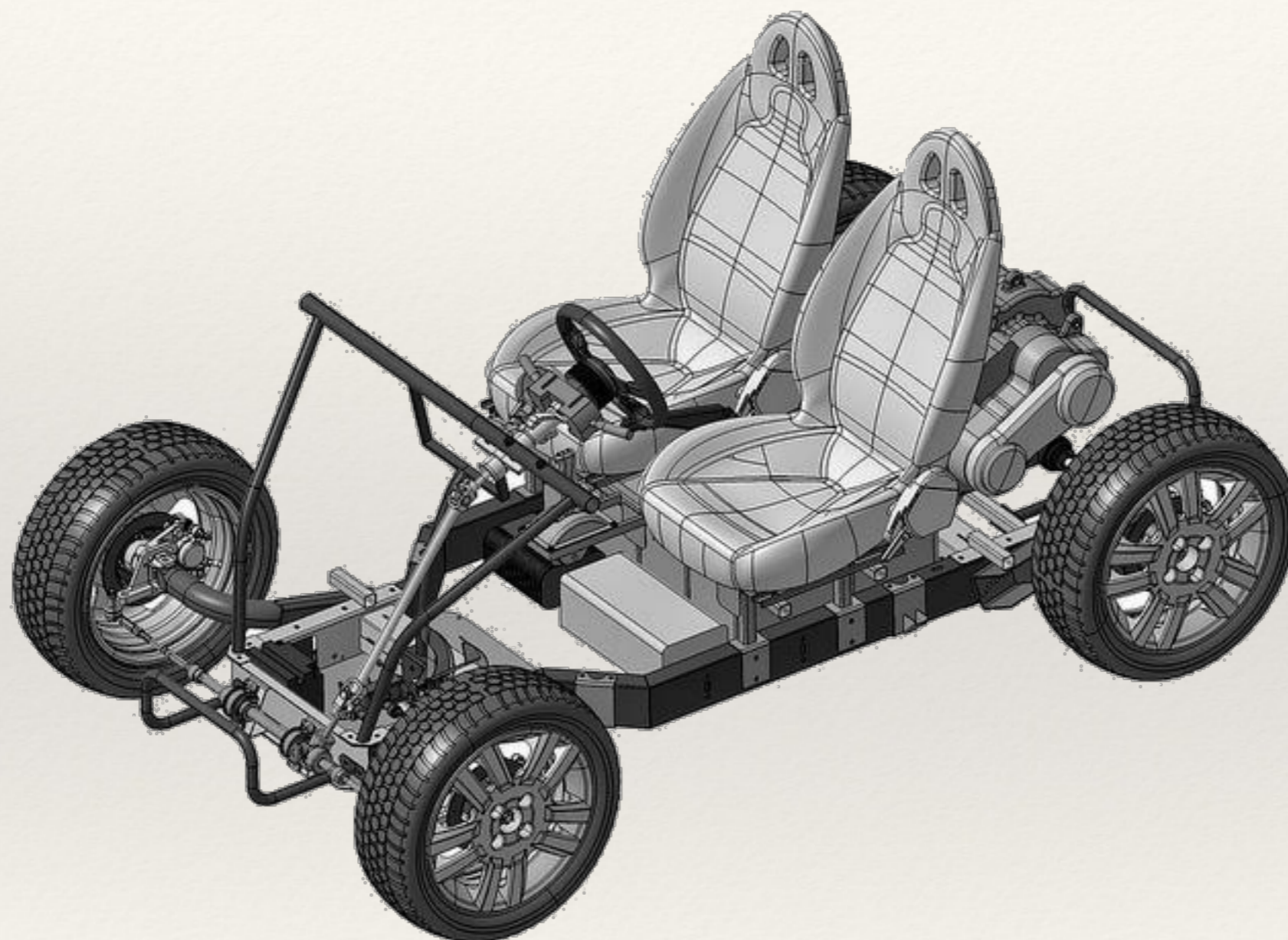


# Evolving

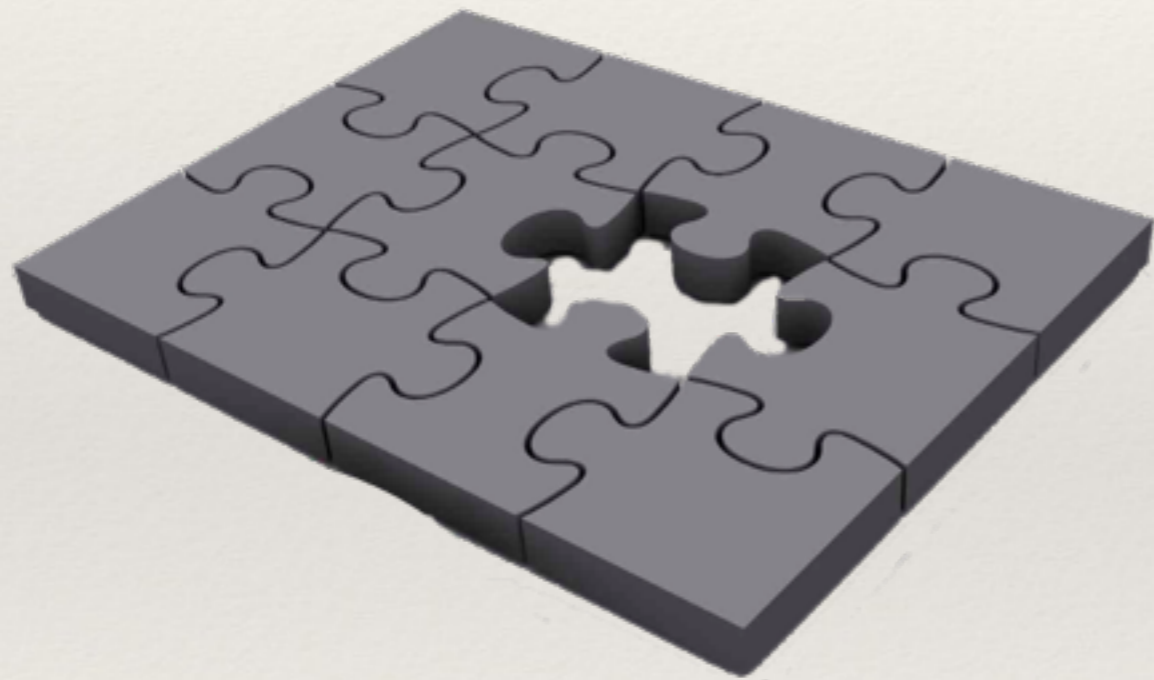




# Incomplete



# Incomplete and Evolving



Incompleteness



Evolution



# Verifying Incomplete and Evolving Specifications

Requirements



Specification



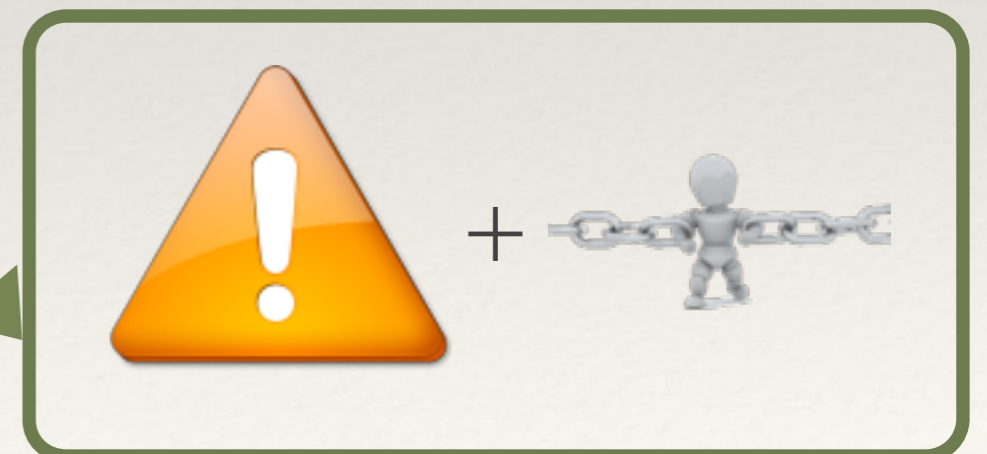
Verification



yes



no + count



maybe

constraints

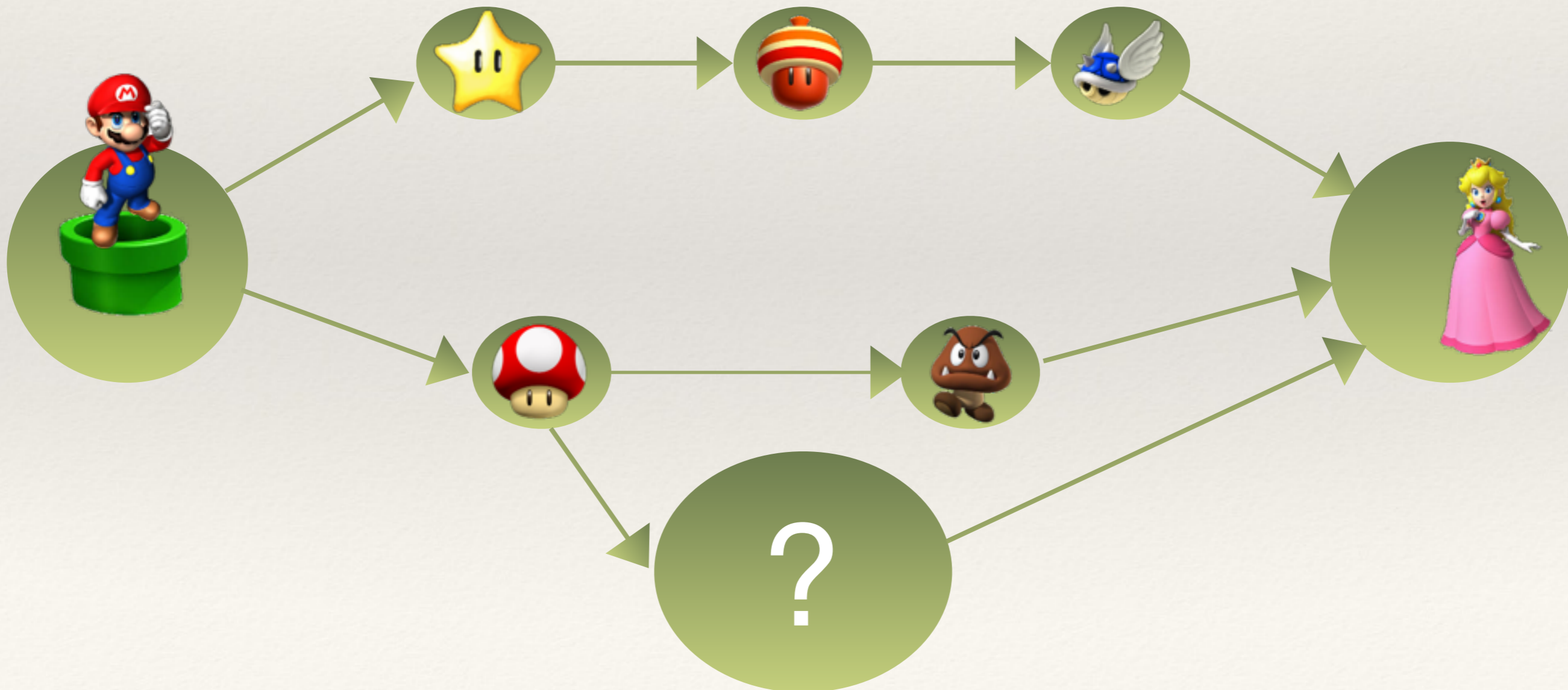
# Constraints

Constraints are conditions (necessary and sufficient) that, if satisfied, guarantee that  $\phi$  holds for  $M$ .



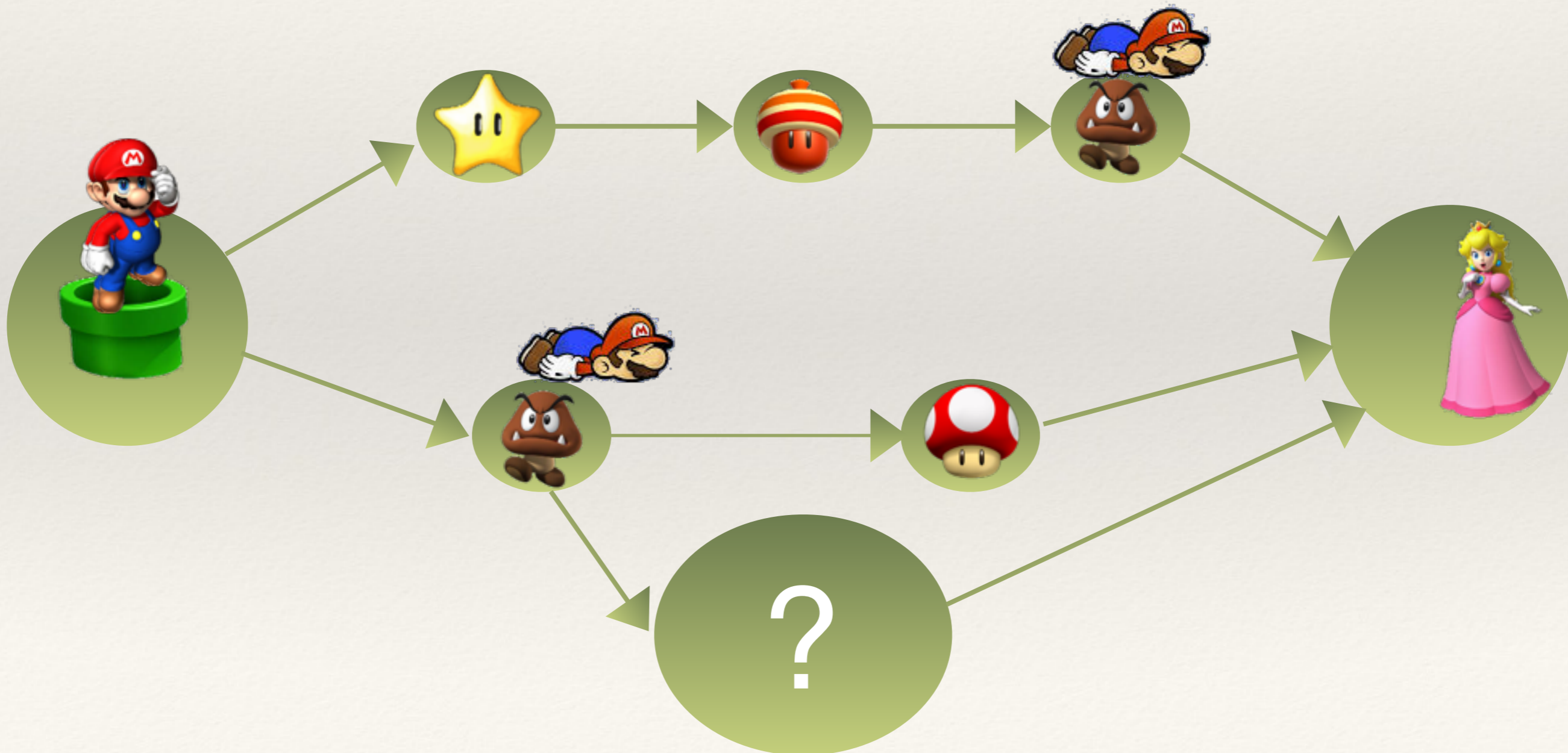
# The idea in practice

$E(\neg(\text{👾})) \cup (\text{👸})$



# The idea in practice

$E(\neg(\text{👾})) \cup (\text{👸})$

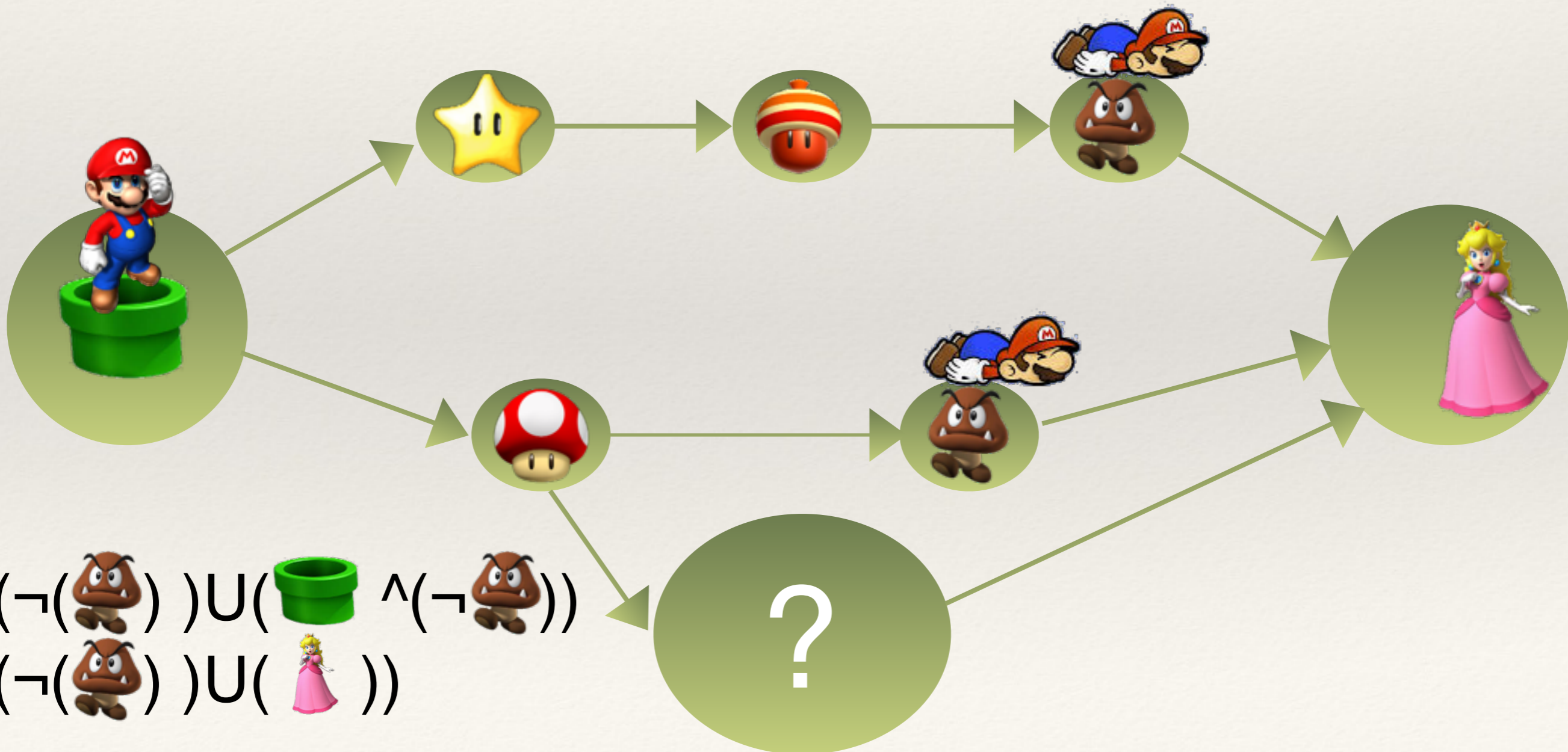




# The idea in practice

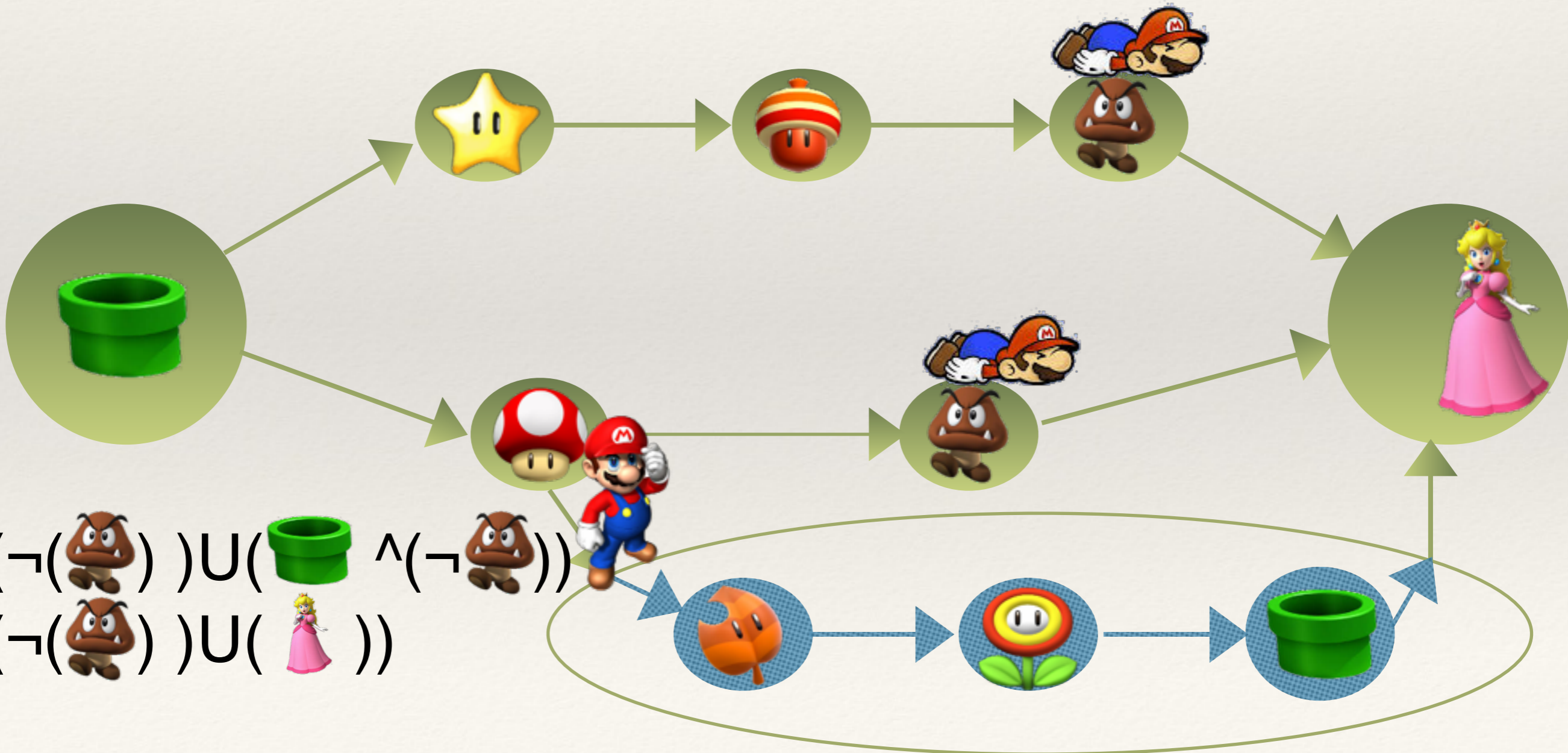


$$E(\neg(\text{👾})) \cup (\text{👸})$$



# The idea in practice

$E(\neg(\text{👾})) \vee (\text{👸})$

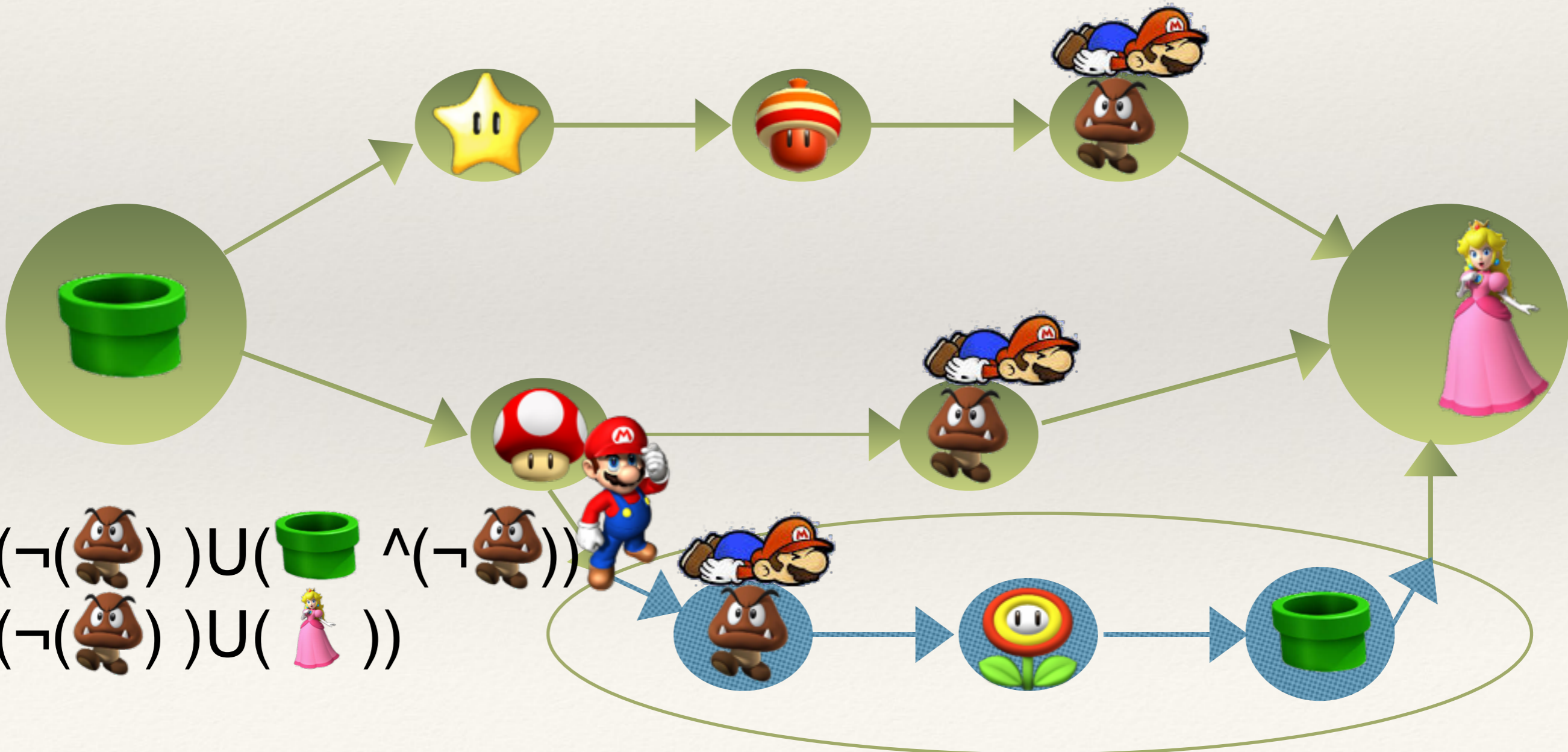




# The idea in practice



$$E(\neg(\text{👾})) \vee U(\text{👸})$$



$$E(\neg(\text{👾})) \vee U(\text{👾} \wedge (\neg(\text{👾})))$$
$$E(\neg(\text{👾})) \vee U(\text{👸})$$

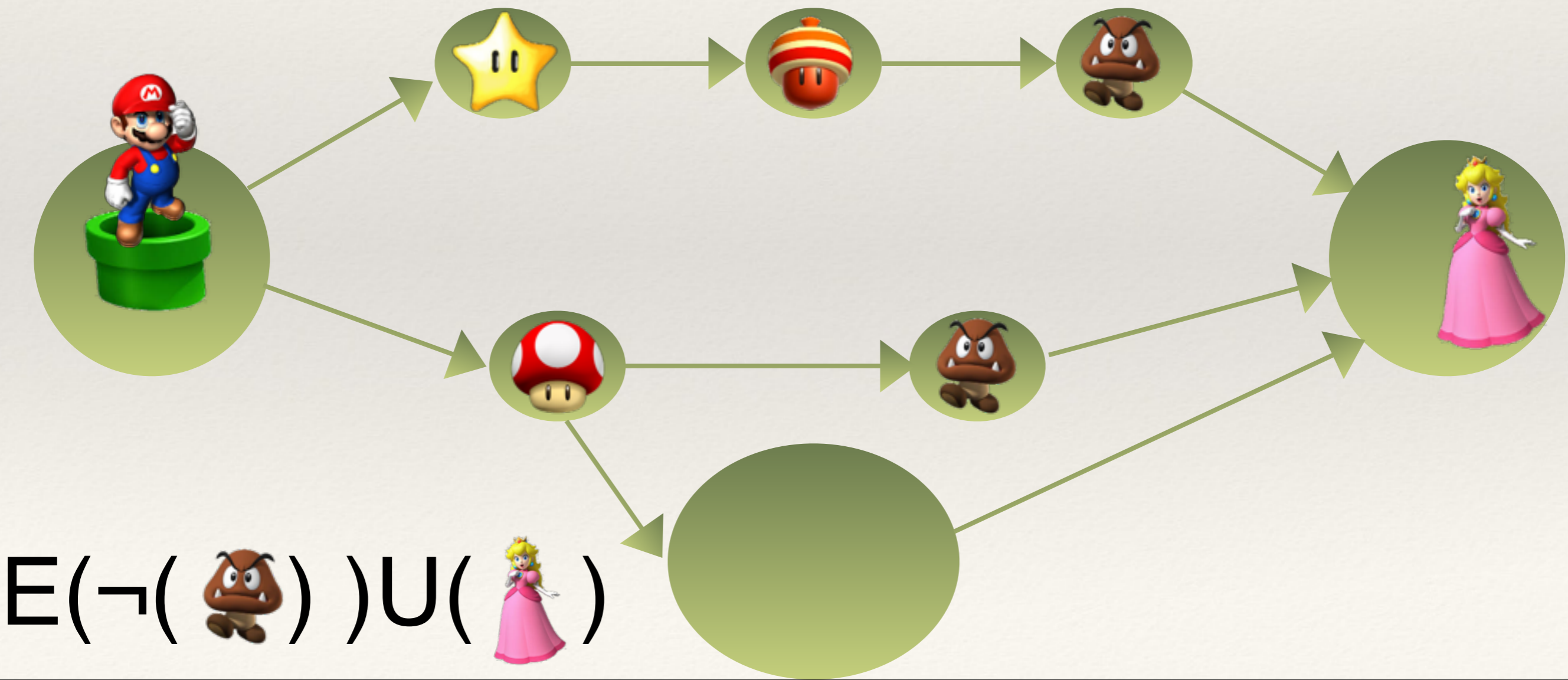
# Consideration of our Approach

The **efficiency** of the approach  
strongly depends  
on the number and on the type  
of the **constraints** generated



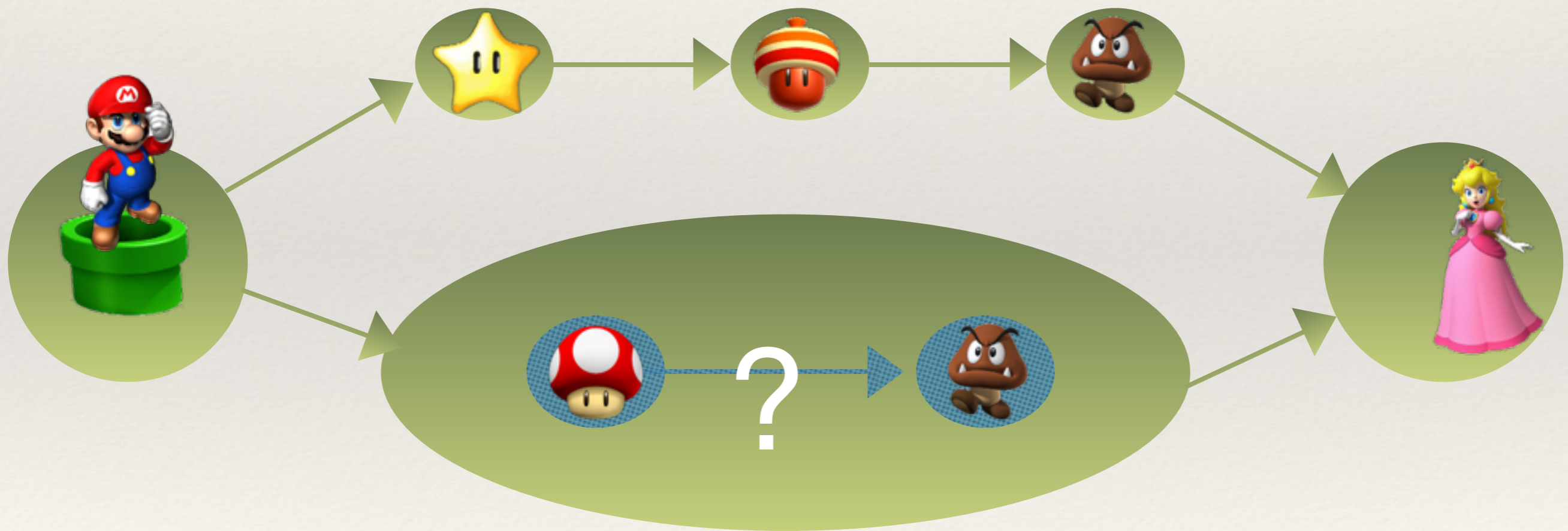
# Consideration of our Approach

*(1) it is possible to check whether an initial, incomplete and high level description of the system (such as the one produced during earlier stages of the software design) satisfies a given property.*



# Consideration of our Approach

(2) *the verification of incomplete models supports the replacement of complex parts of the design with incomplete parts to reduce the verification time*

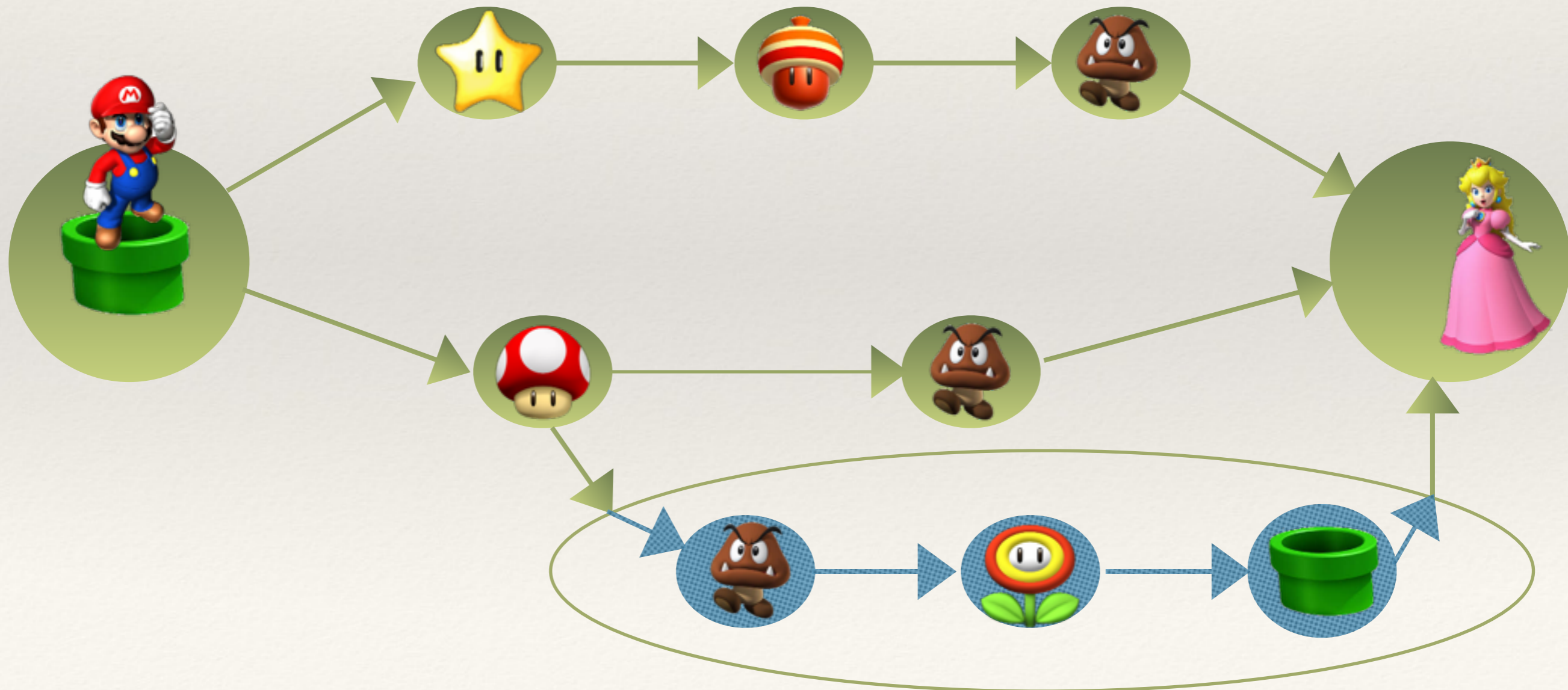


$$E(\neg(\text{Koopa})) \cup (\text{Peach})$$



# Consideration of our Approach

*(3) in the adaptive system case, the verification of incomplete models allows to understand whether the system satisfies its requirements when the different components are removed (plugged) into the running system*



# Formalisms

| <b>Specification</b>                          | <b>Requirements</b>              |
|---|----------------------------------|
| Labeled Transition Systems (LTS) <sup>1</sup> | Computation Tree Logic (CTL)     |
| Buchi automata                                | Linear-Time Temporal Logic (LTL) |
| Statechart <sup>2,3</sup>                     | Computation Tree Logic (CTL)     |
| Markov Chains                                 | Probabilistic CTL                |
| Ambient Calculus                              | Ambient Logic                    |

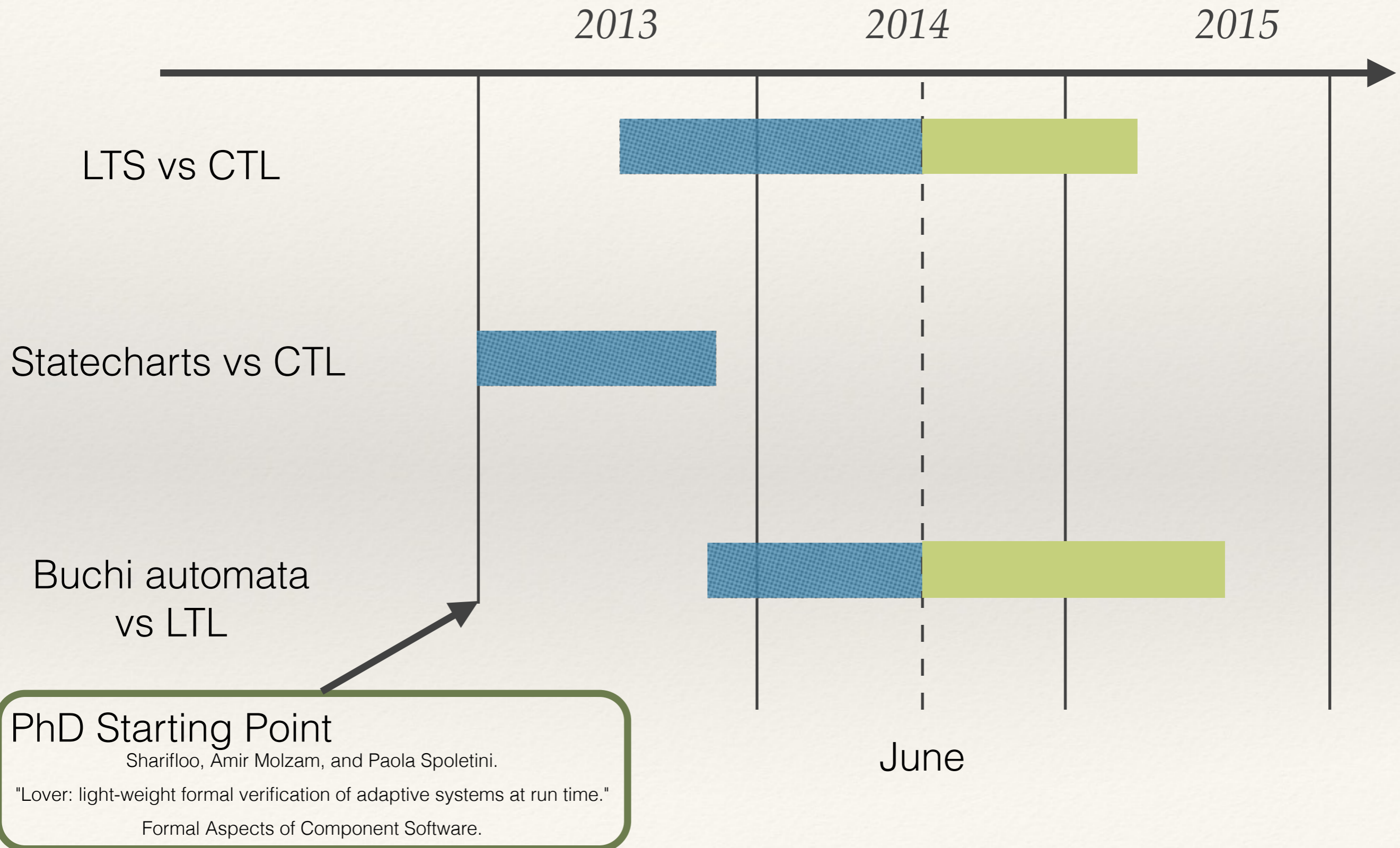
<sup>1</sup> Sharifloo, Amir Molzam, and Paola Spoletini. "Lover: light-weight formal verification of adaptive systems at run time." Formal Aspects of Component Software.

<sup>2</sup> Ghezzi, C., Menghi, C., Sharifloo, A., and Spoletini, P. "On Requirements Verification for Model Refinements." International Requirements Engineering Conference.

<sup>3</sup> Ghezzi, C., Menghi, C., Sharifloo, A., and Spoletini, P. "On requirement verification for evolving Statecharts specifications." Requirements Engineering Journal



# Research Progress



# Thanks to



Carlo Ghezzi



Paola Spoletini



Amir Molzam Sharifloo



Questions ?